

RSA

CS 70 Discussion 4A

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Note: These slides are unofficial course materials. Please use the notes as the only single source of truth.

Problem 1: RSA Intro

(a) Fermat's little theorem: For any prime p

$$a^{p-1} \equiv 1 \pmod{p}$$

$$a^{30} \equiv 1 \pmod{31}$$

FLT allows us to quickly reduce exponents! (But only works for primes)

Euler Totient theorem: For any integer n

$$a^{\phi(n)} \equiv 1 \pmod{n}$$

Where $\phi(n)$ is the number of integers coprime to n !

There is a nice formula for $\phi(n)$ (HW Problem 2)

Problem 1: RSA Intro

(bi) Let $x = 141^{161}$, we want $x \pmod{11 \cdot 17}$

Strategy 1: Use Euler Totient

Strategy 2: Take $\pmod{11}$ and $\pmod{17}$ separately and use CRT to get $\pmod{11 \cdot 17}$

$$x = 141^{161} \equiv 9^{161} \pmod{11} \quad 9^{10} \equiv 1 \pmod{11}$$

$$\equiv 9^{16 \cdot 10 + 1} \pmod{11}$$

$$\equiv 9^{16 \cdot 10} \cdot 9^1 \pmod{11}$$

$$\equiv 1 \cdot 9^1 \equiv 9 \pmod{11}$$

$$x = 141^{161} \equiv 5^{161} \equiv 5 \pmod{17}$$

Problem 1: RSA Intro

(bi) Now apply CRT!

$$x \equiv 9 \pmod{11}$$

$$x \equiv 5 \pmod{17}$$

$$x = 9a + 5b$$

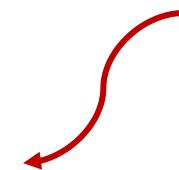
Step 1: “Basis” solutions

$$a \equiv 1 \pmod{11}$$

$$a \equiv 0 \pmod{17}$$

$$b \equiv 0 \pmod{11}$$

$$b \equiv 1 \pmod{17}$$



Step 2: Solve for a and b

$$a = 34$$

$$b = 154$$

Solution:

$$x = 34(9) + 5(154) = 1076 \equiv 141 \pmod{187}$$

Problem 1: RSA Intro

(bi) A (slightly) more clever way

$$x = 141^{161} \equiv 141 \pmod{11}$$

$$x = 141^{161} \equiv 141 \pmod{17}$$

So

$$x - 141 \equiv 0 \pmod{11}$$

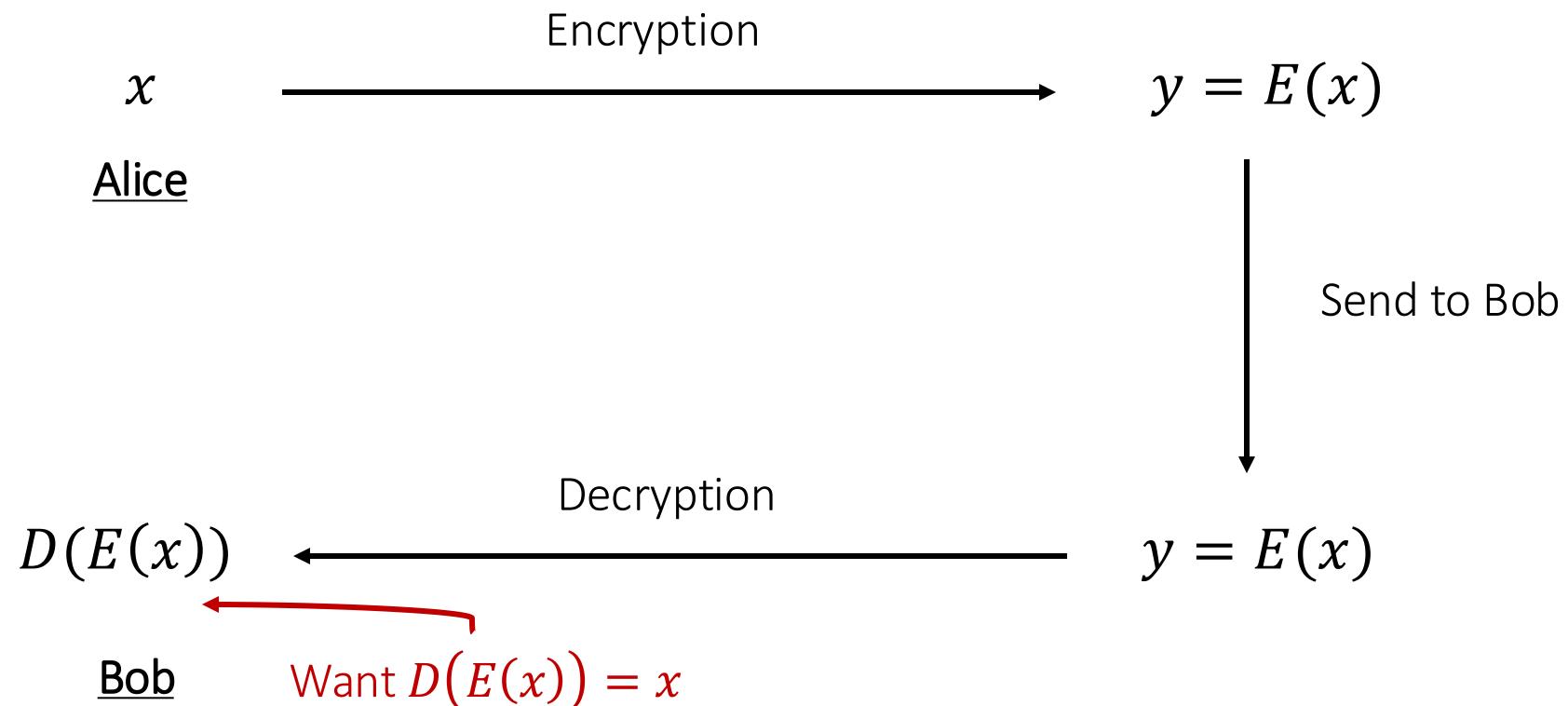
$$x - 141 \equiv 0 \pmod{17}$$

So

$$x - 141 \equiv 0 \pmod{11 \cdot 17}$$

Problem 1: RSA Intro

Alice want to send message x to Bob



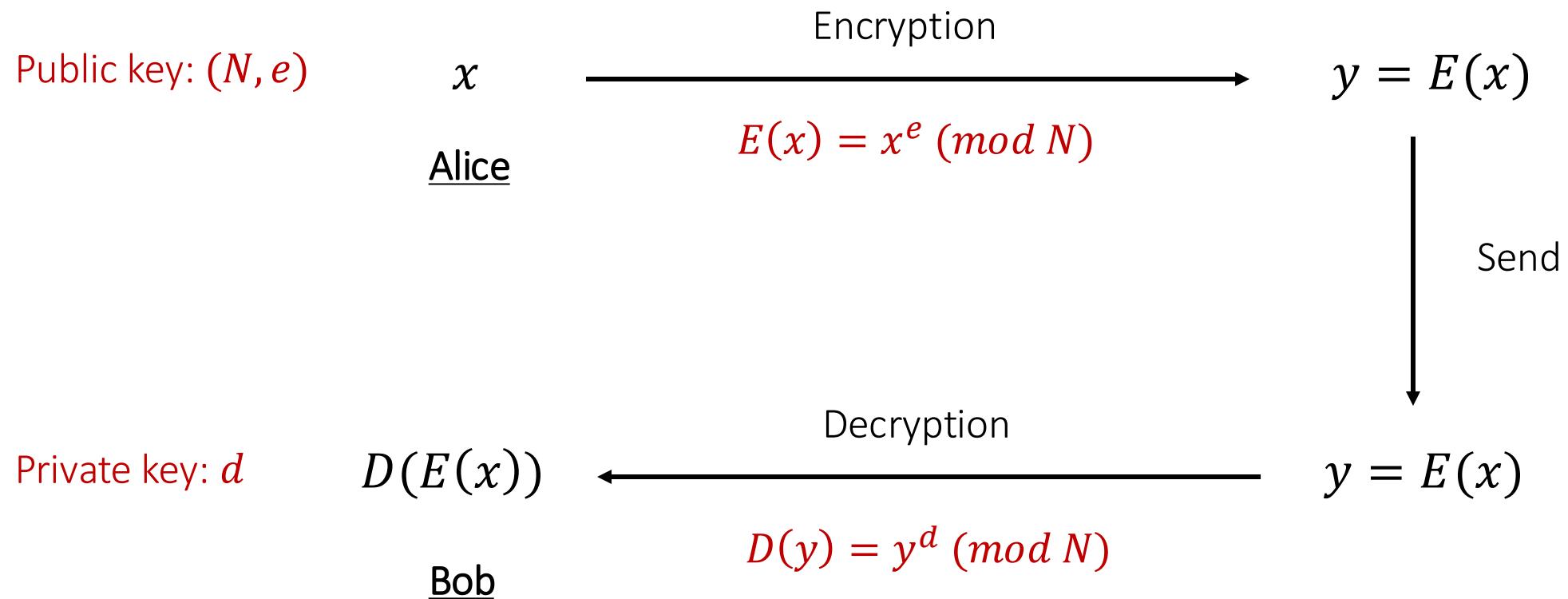
Problem 1: RSA Intro

Setting up RSA

- Step 1: Choose (large) primes p, q
- Step 2: Define $N = pq$
- Step 3: Choose e coprime to $(p - 1)(q - 1)$ ← Public key: (N, e)
- Step 4: Choose $d \equiv e^{-1} \pmod{(p - 1)(q - 1)}$ ← Private key: d

Problem 1: RSA Intro

Alice want to send message x to Bob



Problem 1: RSA Intro

(bii)

Strategy 3: Can we treat $141^{161} \pmod{187}$ as a RSA encoding/decoding scheme

Is this e, d , or ed ?

$$141^{161} \pmod{187}$$

Message: $x = 141$ $N = 187, p = 11, q = 17$

Note that $161 = 7 \cdot 23$

And $7 \cdot 23 \equiv 1 \pmod{160}$

e d

Problem 2: RSA Warm-Up

(a, b, c)

Setting up RSA

- Step 1: Choose (large) primes p, q
- Step 2: Define $N = pq$
- Step 3: Choose e coprime to $(p - 1)(q - 1)$
- Step 4: Choose $d \equiv e^{-1} \pmod{(p - 1)(q - 1)}$

e cannot be even!



$$p = 5, q = 17$$

$$N = pq = 85$$

$$e = 3$$

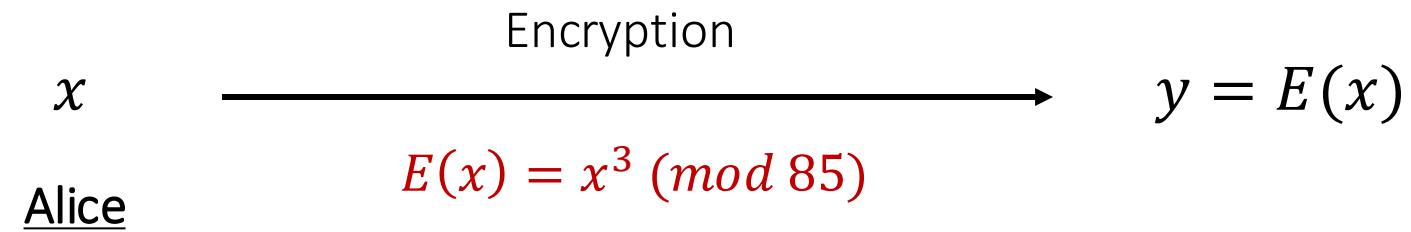
$$d = 3^{-1} \pmod{64} = 43$$

Public key: $(N, e) = (85, 3)$

Private key: $d = 15$

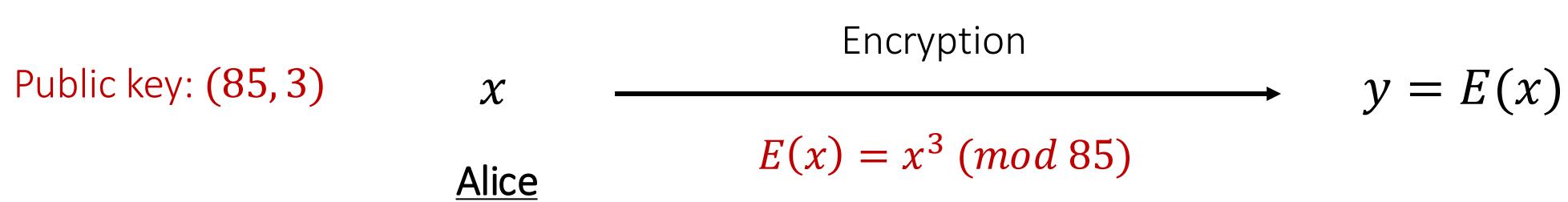
Problem 2: RSA Warm-Up

(d) Alice send message 10 to Bob



Problem 2: RSA Warm-Up

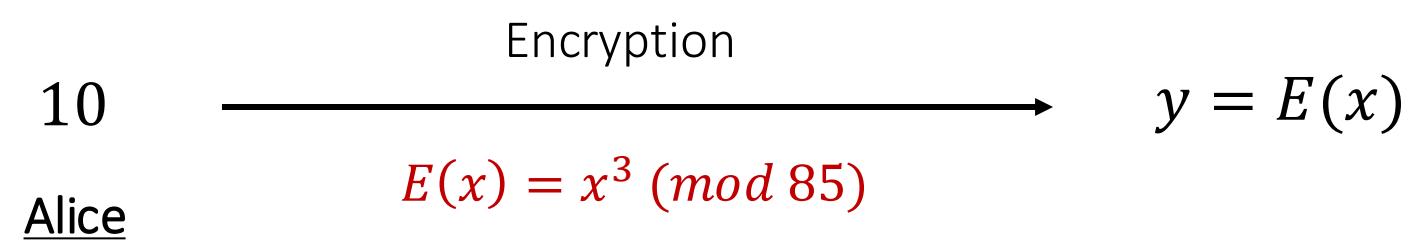
(d) Alice send message 10 to Bob



Problem 2: RSA Warm-Up

(d) Alice send message 10 to Bob

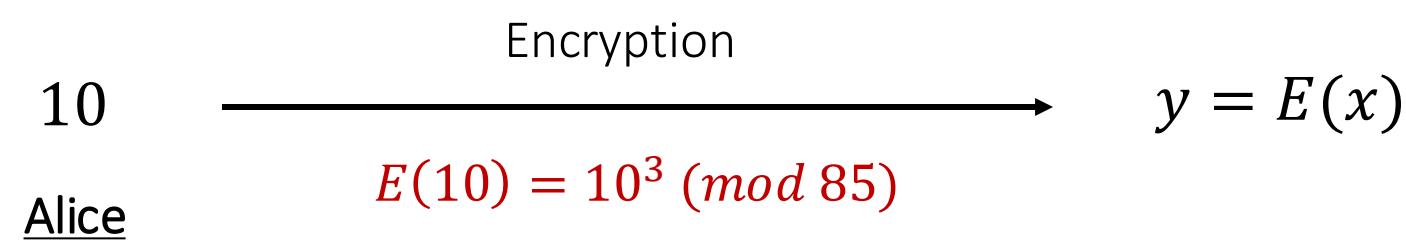
Public key: $(85, 3)$



Problem 2: RSA Warm-Up

(d) Alice send message 10 to Bob

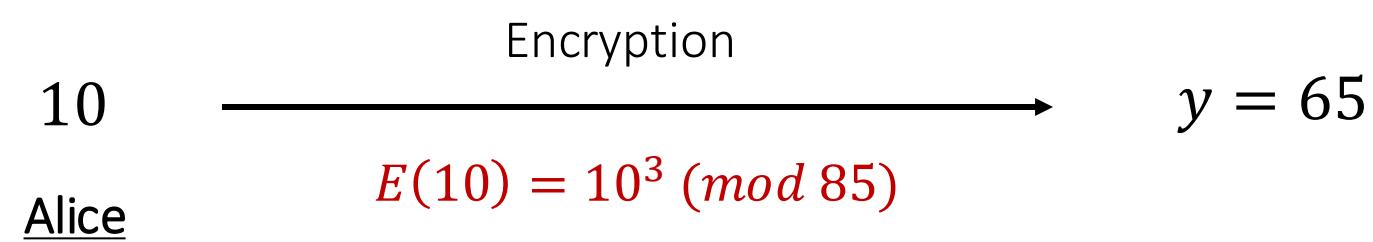
Public key: $(85, 3)$



Problem 2: RSA Warm-Up

(d) Alice send message 10 to Bob

Public key: $(85, 3)$



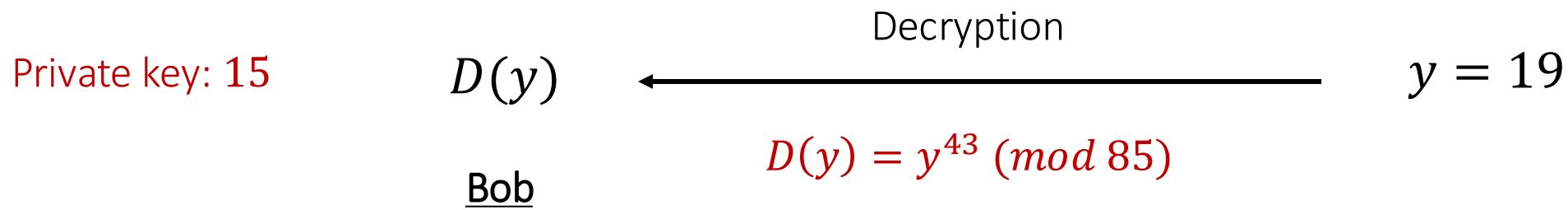
Problem 2: RSA Warm-Up

(e) Bob receives 19 from Alice

$$\begin{array}{ccc} D(y) & \xleftarrow{\text{Decryption}} & y = 19 \\ \underline{\text{Bob}} & & D(y) = y^{43} \pmod{85} \end{array}$$

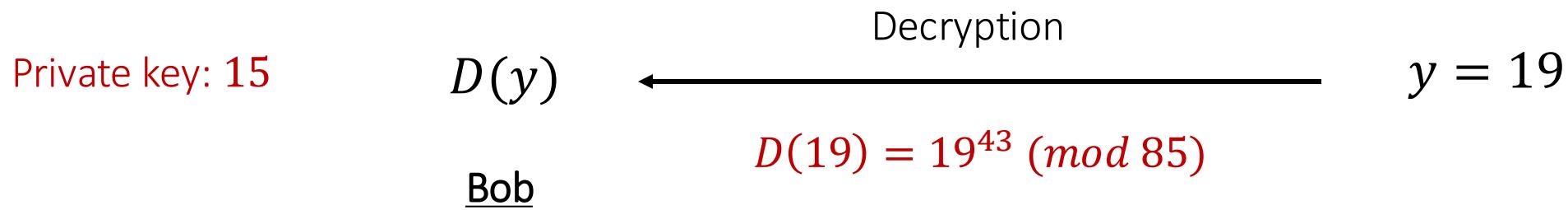
Problem 2: RSA Warm-Up

(e) Bob receives 19 from Alice



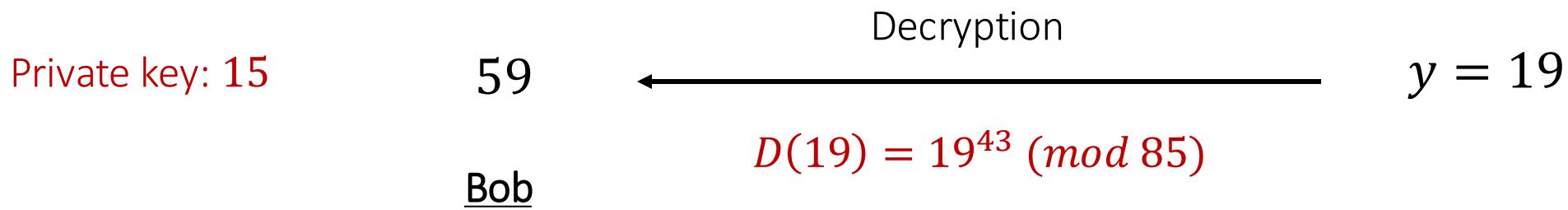
Problem 2: RSA Warm-Up

(e) Bob receives 19 from Alice



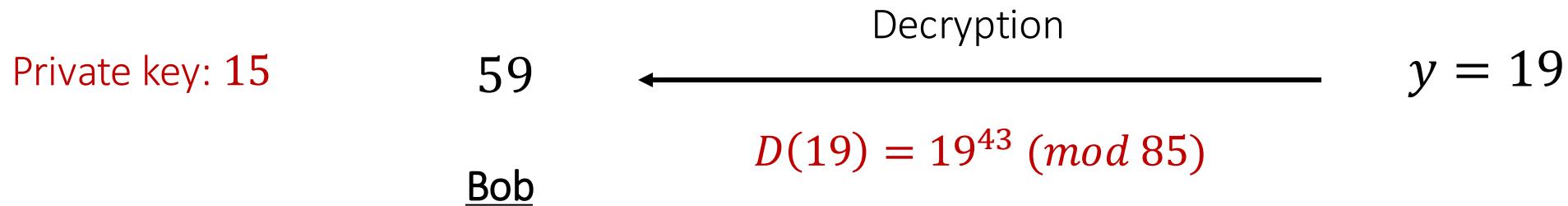
Problem 2: RSA Warm-Up

(e) Bob receives 19 from Alice



Problem 2: RSA Warm-Up

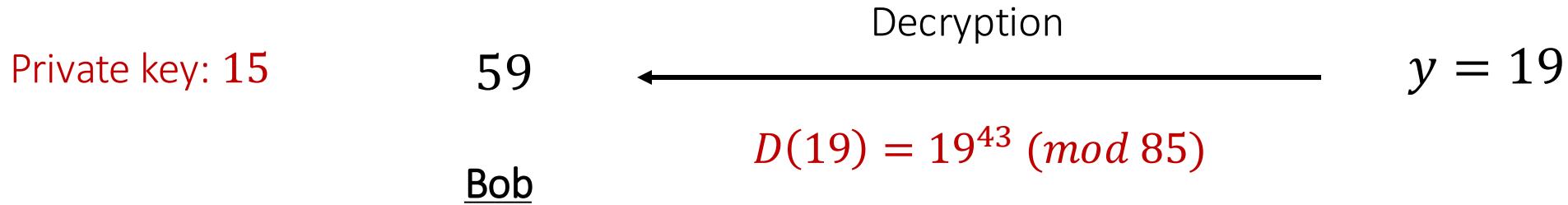
(e) Bob receives 19 from Alice



Step 1: Use FLT to reduce $\text{mod } 5, \text{mod } 17$

Problem 2: RSA Warm-Up

(e) Bob receives 19 from Alice



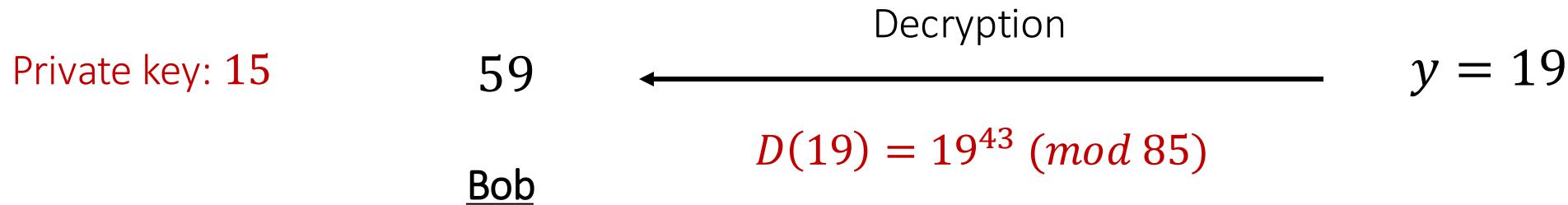
Step 1: Use FLT to reduce $\pmod{5}$, $\pmod{17}$

$$x \equiv 4 \pmod{5}$$

$$x \equiv 8 \pmod{17}$$

Problem 2: RSA Warm-Up

(e) Bob receives 19 from Alice



Step 1: Use FLT to reduce $\pmod{5}$, $\pmod{17}$

$$x \equiv 4 \pmod{5}$$

$$x \equiv 8 \pmod{17}$$

Step 2: Use CRT to find the solution $\pmod{85}$

Problem 2: RSA Warm-Up

(f) Assumptions in RSA

Assumption 1: Factorizing $N = pq$ is hard

- Otherwise we can find the private key!

Assumption 2: Given y it is hard to solve

$$y \equiv x^e \pmod{N}$$

- Otherwise we can just solve for the message x !

Secret word: Chinchilla



Secret word: Meerkat

