

Graph I

CS 70 Discussion 2A

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Note: These slides are unofficial course materials. Please use the notes as the only single source of truth.

Problem 1: Degree Sequence

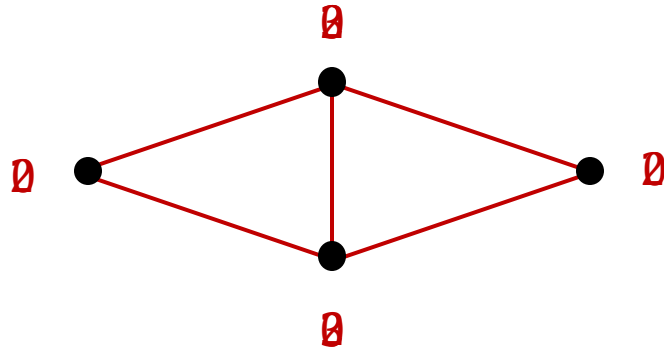
a) (3, 3, 2, 2)

First step: Check using handshaking lemma

Note (Handshaking Lemma):

$$G = (V, E) \text{ is a graph} \implies \sum_{v \in V} \deg(v) = 2|E|$$

Sum of degree is even, so it's "possible" that there exists such graphs



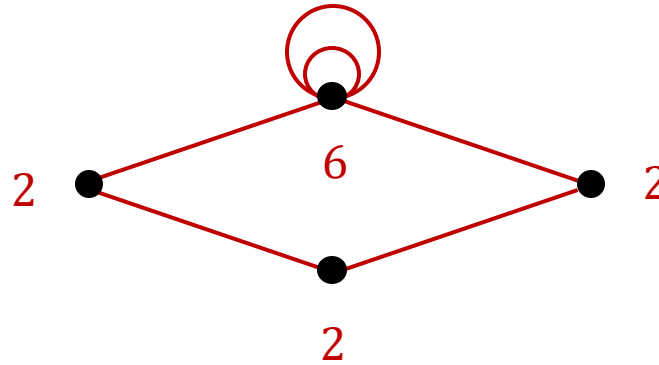
Problem 1: Degree Sequence

b) $(3, 3, 2, 2, 2, 1, 1)$

Sum of degree is odd, so no such graphs exist.

c) $(6, 2, 2, 2)$

Sum of degree is even, so it's possible such graphs exist.

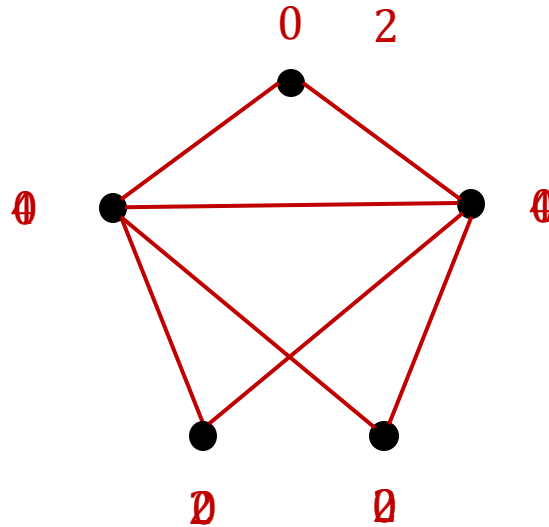


Total number of vertex is 4, so cannot have degree 6 vertex!

Problem 1: Degree Sequence

d) (4, 4, 3, 2, 1)

Sum of degree is even, so it's "possible" such graphs exist.



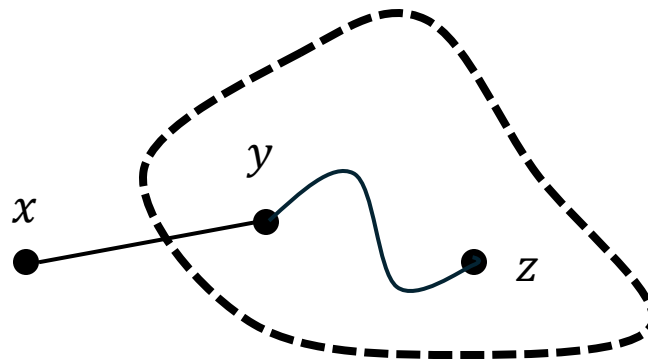
Impossible! The other remaining vertex must have minimum degree of 2!

Problem 2: Build-Up Error

False claim: If every vertex in an undirected graph with $|V| \geq 2$ has degree at least 1, then it is connected

Induction hypothesis: Suppose the statement is true for some $n \geq 2$.

Induction step: We prove the claim for $n + 1$. Consider any undirected graph on n vertices, which every vertex has degree at least 1.



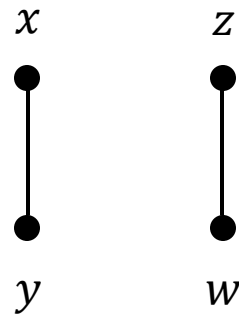
n vertex graph

By IH, it is connected

Problem 2: Build-Up Error

Issue: Not every graph with $n + 1$ vertices can be constructed in such way

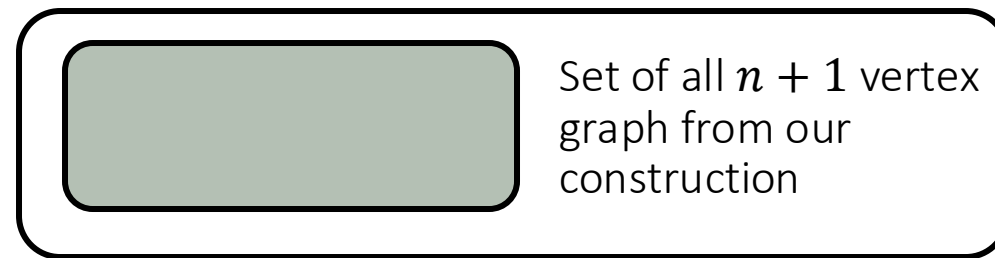
Example:



It's impossible to construct this graph from a 3 vertex graph with min degree 1!

Problem 2: Build-Up Error

Dangerous: start from n -vertex graph and build $n + 1$ -vertex graph



Need to argue that these two sets are the same!

Set of all $n + 1$ vertex graph

Safer: Start from $n + 1$ -vertex graph and remove some vertex to get to n -vertex graph
Then apply induction hypothesis

Problem 3: Eulerian Tour and Walk

a) Is there an Eulerian tour in the given graph?

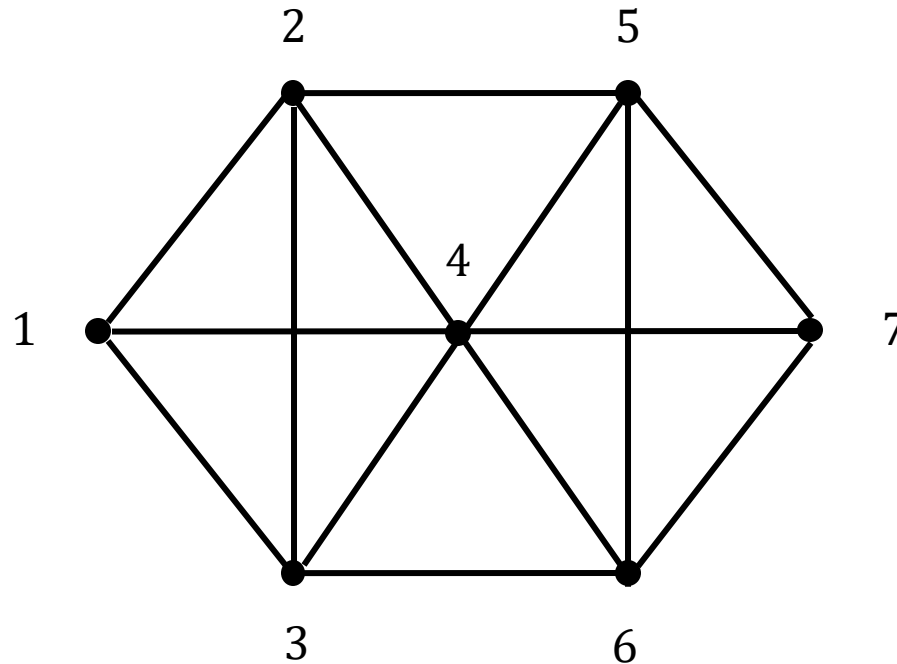
Note (Euler's theorem):

A graph $G = (V, E)$ has an Eulerian tour $\iff G$ is connected and even degree

So no, G is not even degree

Problem 3: Eulerian Tour and Walk

b) Is there an Eulerian walk in the graph below?



Yes: $1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 1 \rightarrow 3 \rightarrow 6 \rightarrow 4 \rightarrow 5 \rightarrow 2 \rightarrow 4 \rightarrow 7 \rightarrow 5 \rightarrow 6 \rightarrow 7$

Problem 3: Eulerian Tour and Walk

c) What is the condition that there is an Eulerian walk in an undirected graph?

A graph G has an Eulerian walk $\iff G$ is connected and has ≤ 2 odd degree vertex

\implies):

Suppose G has an Eulerian walk

\implies Connected, and ...

\implies Only the last two vertex is odd degree

Problem 3: Eulerian Tour and Walk

\Leftarrow):

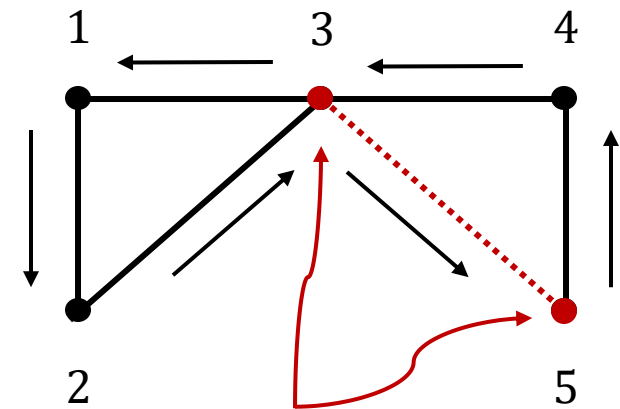
Suppose G is connected and has ≤ 2 odd degree vertex

Connect the two odd degree vertex

$\Rightarrow G$ is now connected and even degree

$\Rightarrow G$ has an Eulerian tour

Removing the added edge gives an Eulerian walk!



Becomes start and end points!

Problem 4: Coloring Trees

a) Prove that all trees with at least 2 vertices have at least two leaves. Recall that a leaf is defined as a node in a tree with degree exactly 1.

Proof 1 (Handshaking Lemma).

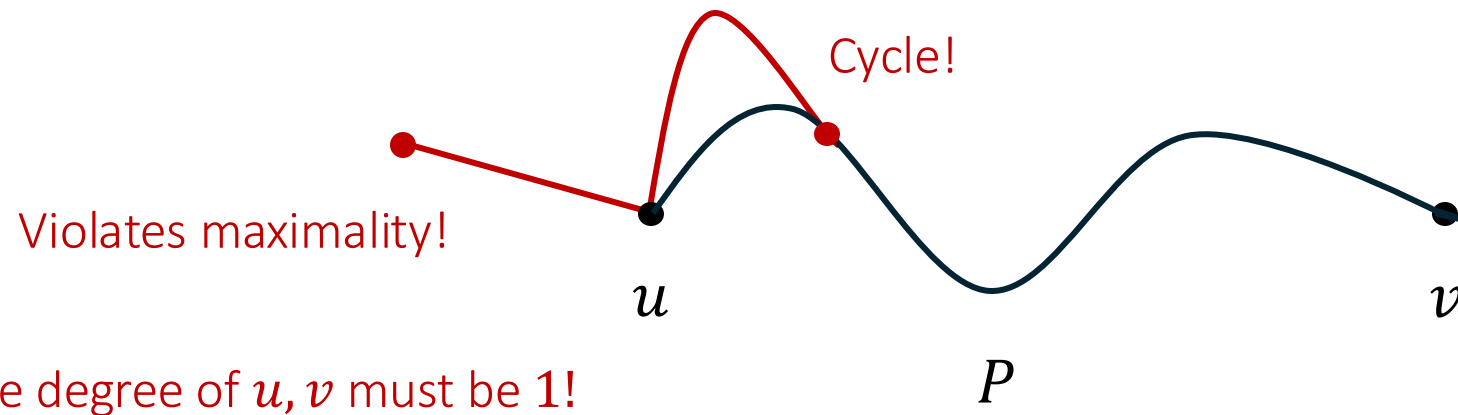
$$\begin{aligned} 2(n - 1) &= \sum_{v \in V} \deg(v) = \sum_{v \in L} \overset{= 1}{\deg(v)} + \sum_{v \in V \setminus L} \overset{\geq 2}{\deg(v)} \\ &\geq \sum_{v \in L} 1 + \sum_{v \in V \setminus L} 2 = |L| + 2(n - |L|) \\ &= 2n - |L| \end{aligned}$$

Problem 4: Coloring Trees

a) Prove that all trees with at least 2 vertices have at least two leaves. Recall that a leaf is defined as a node in a tree with degree exactly 1.

Proof 2 (Maximal path).

Consider the maximal path P in the tree, with endpoints u, v .



The degree of u, v must be 1!

Problem 4: Coloring Trees

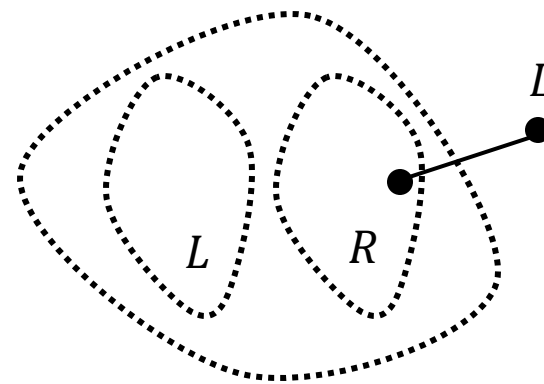
b) Prove that all trees with at least 2 vertices are bipartite.

Proof 1 (Induction).

Induction hypothesis: Suppose the statement is true for trees with n vertex

Induction step: Consider any tree with $n + 1$ vertex

- Remove a leaf
- Apply IH on the n vertex tree
- Reassign the last vertex based on where it's connected



$n + 1$ vertex tree

Problem 4: Coloring Trees

b) Prove that all trees with at least 2 vertices are bipartite.

Proof 2 (BFS).

Use BFS to color the vertices!

